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Reniers, Dennie; Telea, Alexandru

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Skeleton-based Hierarchical Shape Segmentation

Dennie Reniers, Alexandru Telea
Department of Mathematics and Computer Science
Eindhoven University of Technology
PO Box 513, 5600 MB Eindhoven, The Netherlands
d.reniers@tue.nl, alext@win.tue.nl

Abstract

We present an effective framework for segmenting 3D shapes into meaningful components using the curve skeleton. Our algorithm identifies a number of critical points on the curve skeleton, either fully automatically as the junctions of the curve skeleton, or based on user input. We use these points to construct a partitioning of the object surface using geodesics. Because it is based on the curve skeleton, our segmentation intrinsically reflects the shape symmetry and topology. By using geodesics we obtain segments that have smooth, minimally twisting borders. Finally, we present a hierarchical segmentation of shapes which reflects the hierarchical structure of the curve skeleton. We describe a voxel-based implementation of our method which is robust and noise resistant, computationally efficient, able to handle shapes of complex topology, and which delivers level-of-detail segmentations. We demonstrate the framework on various real-world 3D shapes.

1. Introduction

Shape segmentation is the task of decomposing a 3D shape into its meaningful components. In this context, meaningful components are those that a human being intuitively perceives as distinct, logical parts of the shape. Segmentations are useful in shape analysis, shape matching, medical imaging, collision detection, and other geometric processing methods employing divide-and-conquer strategies.

In defining what characterizes meaningful components, several directions have been pursued in the past. One of these is to use curve skeletons. The curve skeleton is a compact shape descriptor, much like a stick-figure representation. The curve skeleton of a 3D shape is a connected 1D structure that is centered within the shape, reflects its circular symmetries, and efficiently captures its topological and geometrical properties [4]. The structure of the curve

skeleton, consisting of branches and junctions, reflects the hierarchy of the meaningful components. The branches are associated with the components, whereas the junctions reflect the relationship between components.

In this paper we present a new framework for hierarchical shape segmentation using the curve skeleton. The curve skeleton is formally defined in terms of geodesics on the object's surface [19]. For each curve skeleton point these geodesics divide the boundary into a set of connected components, providing a natural skeleton-to-boundary mapping that is intrinsic to the definition of the curve skeleton. Being geodesics, segment borders are smooth and minimally twisting. After computing the meaningful components, we combine them into a hierarchical level-of-detail segmentation. Although we use the curve skeleton junctions to provide the meaningful components, other points can be used as well, depending on the target application. Besides the segmentation framework, this paper contributes a robust algorithm for computing the curve skeleton of voxelized objects which enhances the method in [19], and a method to robustly detect junctions.

Figure 1 provides an overview of our approach, consisting of four stages. First, the curve skeleton is computed from the voxelized input shape. Second, the *critical points* are chosen either automatically as the junctions of the curve skeleton, or manually by the user. Third, we compute the component sets of the critical points. Finally, a level-of-detail segmentation is created from the component sets.

The outline of this paper is as follows. In the next section we review related work. In Section 3 we present some preliminary definitions. In Section 4 we provide the details of the curve skeleton definition and computation. In Section 5 we elaborate on the computation of the component sets, and how based on them junctions can be detected robustly. Section 6 deals with creating the hierarchical segmentation from the component sets. In Section 7 we present the results of the algorithm. Section 8 presents a discussion and we conclude in Section 9.

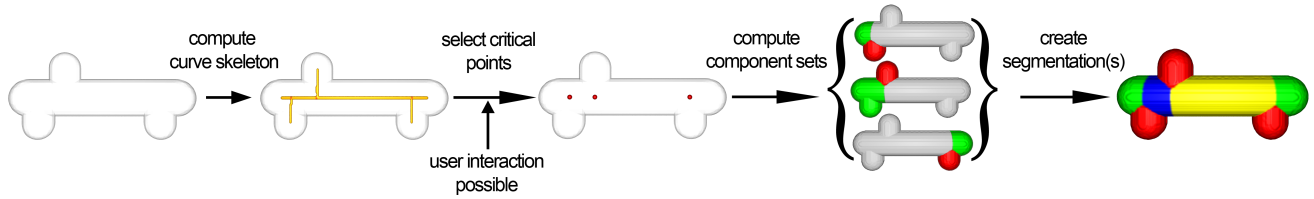


Figure 1. Overview of our segmentation framework.

2. Related Work

Shapes can be segmented by considering their boundary or interior. Approaches that are based on the boundary usually define segment borders at object-surface concavities. Katz and Tal [11] use fuzzy clustering on geodesic and angular distances between surface elements to obtain a hierarchical mesh decomposition with non-jaggy borders. They show that their segmentation can be used to compute a (control) curve skeleton. In [10], pose-invariant segmentations are produced by extracting feature points and cores. In [3], a fuzzy clustering on quasi-flat surface features separated by curvature extrema is used to obtain a multiscale segmentation of point sets. In [14], an automatic scissoring scheme is proposed based on 3D snakes.

Segmentation methods using the curve skeleton consider the shape's interior. Our approach falls into this category. To obtain meaningful components, these methods require a mapping from the curve skeleton to the object boundary. Various mappings have been proposed. In [15] for example, a combination of planar cross-sections and space sweeping is used. The approach in [5] uses force-following of boundary particles onto the curve skeleton. A comparative study of some of the latest segmentation methods can be found in [2].

Our segmentation approach is similar to that of Li et al. [15], in which the segmentation is also based on curve skeletons. However, there are some important differences. Our curve skeleton algorithm has a more fundamental underpinning (see [19] for details), is more noise resistant, and is connected by default. Li et al. use a planar cross-section sweeping along the curve skeleton between critical points to obtain components. The definition of our curve skeleton is such that it provides a natural skeleton-to-boundary mapping. Instead of taking the planar cross-section as borders between components, we use the actual curve skeleton definition and take the shortest geodesics between feature points as borders. This makes for more natural borders between segments (compare e.g. the Hand object in Fig. 11 with [15, Fig.13]). Furthermore, our framework provides a hierarchical segmentation.

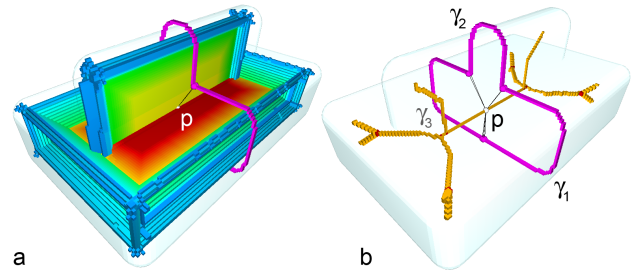


Figure 2. The surface skeleton with rainbow color map indicating geodesic length (blue is short, red is long) (a). The curve skeleton \mathcal{C} , a selected point $p \in \mathcal{C}$, and $\Gamma(p) = \{\gamma_1, \gamma_2, \gamma_3\}$ (b).

3. Definitions

The surface skeleton \mathcal{S} of a 3D object Ω with boundary $\partial\Omega$ is defined as those points in Ω having at least two boundary points at minimum distance:

$$\mathcal{S} = \{p \in \Omega \mid \exists a, b \in \partial\Omega, a \neq b, \text{dist}(p, a) = \text{dist}(p, b) = \min_{k \in \partial\Omega} \text{dist}(p, k)\}, \quad (1)$$

where dist is the Euclidean distance in \mathbb{R}^3 . Points a, b are the points at minimum distance of p and are called the *feature points* of p . Let $F : \Omega \rightarrow \mathcal{P}(\partial\Omega)$ be the *feature transform* which assigns to each object point its set of feature points, where \mathcal{P} denotes the power set. The surface skeleton consists of 2D manifolds, called *sheets*, and of 1D curves [8]. The sheets intersect in *sheet-intersection curves*. Points on these intersection curves have at least three feature points. The algorithms we present in this paper all work in binary voxel space [13]. When used in a discrete-space context, Ω denotes the set of foreground, or object voxels.

4. Curve skeleton computation

First, we define the curve skeleton \mathcal{C} for the generic case, namely when the curve skeleton is not (partly) incident with a sheet-intersection curve, so that there are exactly

two feature points for each point $p \in \mathcal{C}$. The curve skeleton is defined as those points having at least two shortest geodesics γ_i, γ_j between their two feature points $F(p)$:

$$p \in \mathcal{C} \Leftrightarrow \exists \text{ two shortest geodesics } \gamma_i \neq \gamma_j \text{ between } F(p) \quad (2)$$

A similar definition was first presented in [7], and in [19]. This definition ensures that the curve skeleton is centered in two ways. First, the curve skeleton is included in the surface skeleton, as the surface skeleton is defined as those points having at least two feature points. Second, it is also centered within the surface skeleton structure, as the two shortest geodesics γ_i, γ_j necessarily have the same length.

In the non-generic case, the curve skeleton is partly incident with a sheet-intersection curve, and there are more than two feature points for a point $p \in \mathcal{C}$. For example, Figure 2a depicts the surface skeleton of a box with a vertical ridge. Figure 2b depicts the curve skeleton of the same object. Point p lies on the sheet-intersection curve of the box's sheet and the ridge's sheet, and has three feature points. Point p clearly is a curve skeleton point, but definition (2) does not detect this. Consequently, this definition cannot be used to detect all curve skeleton points, so we modify it as follows. With each intersecting sheet a feature-point pair a, b is associated and thus a shortest geodesic γ_i between a, b . For example, for three intersecting sheets there are three shortest geodesics (Fig. 2b), not necessarily of the same length. We call the combination of these shortest geodesics for a surface skeleton point $p \in \mathcal{S}$ the *shortest-geodesic set* $\Gamma(p)$:

$$\Gamma(p) = \{\gamma_i\}_i, \quad (3)$$

where γ_i is a shortest geodesic between the feature-point pair $a, b \in F(p)$ associated with a sheet.

Both in the generic and non-generic case, the union of shortest geodesics form a Jordan curve on the object surface $\partial\Omega$. Hence, detecting a curve skeleton point p comes down to detecting whether $\Gamma(p)$ contains a ring, i.e. computing the genus of $\Gamma(p)$. By the genus of Γ , we mean the genus of the surface that is obtained by taking a infinitesimal dilation of Γ on $\partial\Omega$. The new definition of the curve skeleton replaces (2) and becomes:

$$p \in \mathcal{C} \Leftrightarrow \text{genus}(\Gamma(p)) \geq 1 \quad (4)$$

Additionally, the genus of Γ can be used to differentiate between junction and non-junction points, called *regular* points. A junction is a point on the curve skeleton where at least three branches come together. The shortest-geodesic set of such a junction is the union of the Jordan curves of the neighboring regular points. Hence, if the genus of $\Gamma(p)$ is larger than 1, p is a junction:

$$p \in \text{junctions}(\mathcal{C}) \Leftrightarrow \text{genus}(\Gamma(p)) \geq 2 \quad (5)$$

As mentioned, a similar curve skeleton definition was first presented by Dey and Sun [7], namely as the set of singular points of the so-called medial geodesic function, which assigns to each surface skeleton point the length of the geodesic between its two feature points, and they show that this curve skeleton is homotopy equivalent to the original shape. In contrast, we detect the curve skeleton by performing a topological analysis of the shortest geodesic set. This definition allows us to compute the curve skeleton without first detecting the surface skeleton, and readily allows junction detection. At an implementation level, whereas [7] uses polygonal representations, our implementation of the curve skeleton detection is voxel-based, as detailed in the next section.

4.1. Algorithm

Based on the above definitions, we now present our algorithm for computing the curve skeleton of a voxelized object Ω . The algorithm consists of several steps that can be executed in parallel for each object voxel. The different stages are depicted in Figure 3. In the first step, we compute the feature transform F and extended feature transform \bar{F} of Ω . The *extended feature transform* merges the feature set of each voxel p with that of its first octant 26-neighbors: $\bar{F} = \bigcup_{x,y,z \in \{0,1\}} F(p_x + x, p_y + y, p_z + z)$. The result is that a regular curve skeleton point has two (compact) groups of feature voxels (Fig. 3a), one on each side of the surface skeleton. The extended feature transform solves two problems. First, it solves the well-known problem that in discrete space the feature set of a surface skeleton voxel might contain only one feature voxel, since there might be no voxel exactly in the center of the voxelized boundary [6, 18]. Second, it solves the related problem that on a voxelized object surface, there is typically only one shortest path between the two feature voxels of a curve skeleton voxel, although there are two on the underlying continuous surface.

After computing \bar{F} , we compute the shortest-geodesic set $\Gamma(p)$. Because we perform our computations in discrete voxel space, we compute the shortest geodesics as shortest paths in the *boundary graph* in which the surface voxels are the nodes and their neighborhood relations represent the edges. The shortest paths are represented as 3D chain codes [12] in the boundary graph. For computing a shortest path we use the A* algorithm [9], a modification of Dijkstra's algorithm, with Euclidean distance as the search heuristic. We compute a shortest path between each two feature voxels in $\bar{F}(p)$. Because $\bar{F}(p)$ contains groups of voxels, we will find numerous shortest paths for a regular point having only two feature points. Together, they form a discrete (noisy) Jordan curve on the object surface (Fig. 3b).

After computing the path set $\Gamma(p)$, we determine its

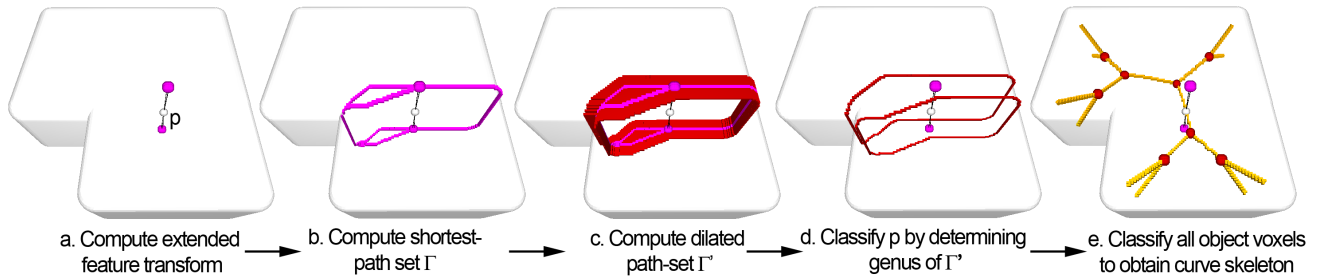


Figure 3. Curve skeleton computation and junction detection.

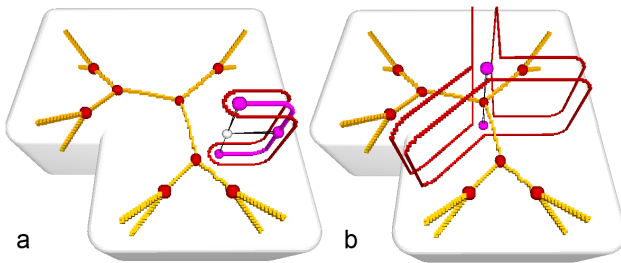


Figure 4. The boundaries of two dilated path-sets.

genus on $\partial\Omega$, as follows. First we dilate $\Gamma(p)$ so that we obtain a surface band $\Gamma'(p)$ centered at the 1D structure $\Gamma(p)$ (Fig. 3c). Conveniently, small holes from the path set that are caused by the discrete nature of the shortest paths (Fig. 3b) are removed by the dilation. The dilation is performed by propagating the voxels in $\Gamma(p)$ a short distance outward using a distance-ordered flood fill. Next, we determine the number of compact boundary pieces that the area $\Gamma'(p)$ has. If two or more boundaries are found, p is concluded to be a curve skeleton voxel (Fig. 3d). When only one boundary is found, the voxel is concluded to be a non-curve skeleton voxel (Figure 4a). Empirical studies on an extensive family of real-world 3D shapes show that a (geodesic) dilation distance of 5.0 is enough to obtain two disjoint boundaries if $\Gamma(p)$ is the discrete representation of a Jordan curve. In principle, when three or more boundaries for $\Gamma'(p)$ are found, p is a junction. In Figure 4b for example, three boundaries are found for the junction voxel. However, due to the discrete nature of the shortest paths, curve skeleton voxels might be incorrectly classified as junctions, i.e., junction detection by analyzing the genus of Γ is conservative. This is problematic if we want to use the junctions in segmentation (Sec. 6). We address this problem in Section 5.4.

The steps described above (Fig. 3a-d) are executed for all object voxels in parallel, resulting in the curve skeleton \mathcal{C} (Fig. 3e). Figure 5 shows four curve skeletons as computed

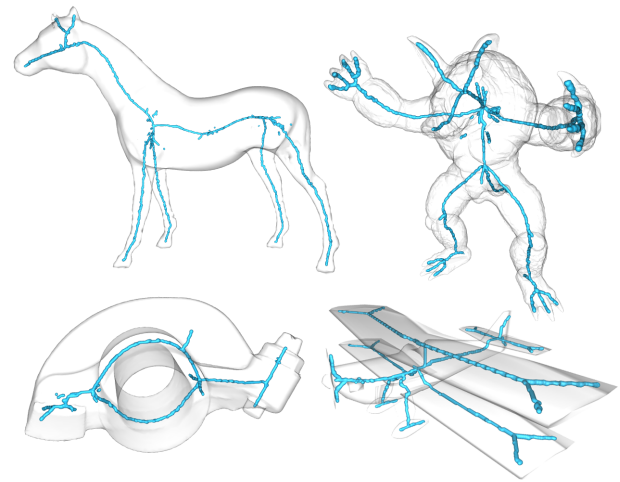


Figure 5. The curve skeleton \mathcal{C} of several objects.

by our algorithm. We stress that these results are obtained without any post-processing, and that our algorithm does not perform any thinning or erosion step. The curve skeletons, and consequently their junctions, are up to 2 voxels thick, due to the extended feature transform. We observe that the curve skeleton is homotopic to the object, and is connected, although some noise is present. In Section 5.3 we present an importance measure for \mathcal{C} that can be used to obtain noise-resistant, simplified skeletons. This measure is based on the component sets that are associated with each point $p \in \mathcal{C}$, as explained in the next section.

5. Component Sets

In the previous section we described the detection of a curve skeleton point p by analyzing the genus of its associated geodesic set $\Gamma(p)$. In this definition of the curve skeleton lies a natural skeleton-to-boundary mapping. For objects of genus 0, the Jordan curve theorem states that the

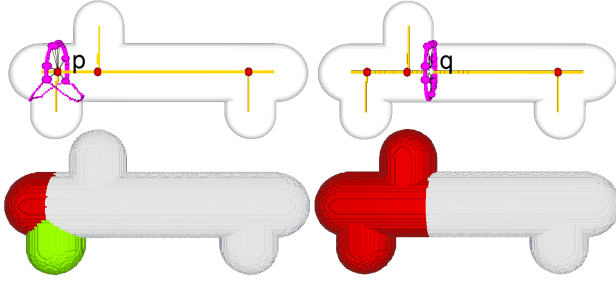


Figure 6. The component sets (bottom) of two selected points p, q (top).

Jordan curve Γ associated with a regular point divides the object surface $\partial\Omega$ into two connected components. For a junction point, Γ is the union of Jordan curves of neighboring regular points, dividing the boundary into multiple components. The *component set* of a point p is denoted $C(p)$. Let the k components in C be ordered by their areas: $\forall_{1 \leq i < k} |C_i| \leq |C_{i+1}|$. The area of a component $C_i(p)$ is determined simply by counting its voxels: the cardinality of the set $C_i(p)$. Although we could use a more sophisticated surface estimator, the cardinality is sufficient for our purposes. In our voxel space representation, the component set $C(p)$ for a voxel p is computed by labeling the connected components in the boundary graph from which the voxels from the path set $\Gamma(p)$ are removed. The labeling is sped up using a simple spatial-subdivision scheme.

Figure 6 shows the component sets for two selected curve skeleton points p, q , where p and q are a junction and regular point respectively. The component sets of the junctions are used to obtain a segmentation in Section 6.

5.1. Topological Properties

It is important to note that the number of components $|C(p)|$ for a point $p \in \mathcal{C}$ is related to whether the shape Ω has tunnels, as this will affect the final segmentation. For shapes without tunnels or holes (genus 0), the curve skeleton has a tree structure. For objects with tunnels (genus ≥ 1), the curve skeleton is a graph that contains a *loop* around each tunnel. Since there is a skeleton-to-boundary mapping through the use of Γ , whether $\Gamma(p)$ divides the boundary into multiple components is equivalent to whether p divides the graph into multiple components. In graph theory, a point p dividing the graph into multiple connected components is called a *cut vertex* [22]. The number of components $|C(p)|$ is related to the genus of $\Gamma(p)$ and the amount of loops $L(p)$ that p is on:

$$|C(p)| = \max(\text{genus}(\Gamma(p)) + 1 - L(p), 1). \quad (6)$$

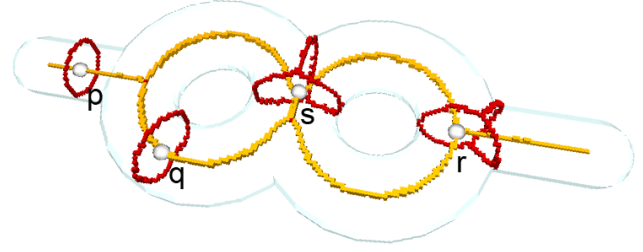


Figure 7. An object with two tunnels and Γ of four selected points.

Figure 7 shows an example shape with two tunnels and four selected points. Points p, q are regular (non-junction) points ($\text{genus}(\Gamma(p)) = 1$, $\text{genus}(\Gamma(q)) = 1$), whereas r, s are junctions of three branches ($\text{genus}(\Gamma(r)) = 2$, $\text{genus}(\Gamma(s)) = 2$). Point p is not on a loop, points q, r are both on one loop, and s is on two loops. Indeed, we verify that the cardinality of the component sets generated by p, q, r, s is 2, 1, 2, 1 respectively. Points p, r are cut vertices in the curve skeleton graph, whereas q, s are not.

5.2. Inclusion Property

An important property of the component sets for the purpose of hierarchical segmentation is that any two components $C_i(p), C_j(q)$ for two points $p, q \in \mathcal{C}$ cannot partly overlap:

$$C_i(p) \cap C_j(q) \in \{\emptyset, C_i(p)\} \quad \text{if } |C_i(p)| \leq |C_j(q)| \quad (7)$$

In order to prove (7) for two regular points p, q , we have to prove that the Jordan curves $\Gamma(p)$ and $\Gamma(q)$ cannot intersect. The proof can be roughly sketched as follows. Take that subset $\mathcal{C}_{pq} \subset \mathcal{C}$ that is between p and q . It is reasonable to assume that the feature points $F(\mathcal{C}_{pq})$ generate two curves L_1, L_2 on $\partial\Omega$, one on each side of the surface skeleton. Each geodesic $\gamma \in \Gamma$ has one endpoint on L_1 and the other endpoint on L_2 . Such a geodesic γ does not intersect L_1 or L_2 in any other point as it would no longer be the shortest geodesic. Now, if a geodesic $\gamma_p \in \Gamma(p)$ intersects $\gamma_q \in \Gamma(q)$, it needs to have (a multiple of) two intersection points. But if γ_p intersects γ_q in two points a, b , the shortest geodesic between a, b would not be unique, which is impossible.

5.3. Importance Measure

The component sets can be used to define an importance measure for computing simplified curve skeletons that are robust to noise, as follows. The *importance measure* $\rho(p)$

for a point $p \in \mathcal{C}$ is defined as the importance of p in representing the original shape. One way of expressing $\rho(p)$ is in terms of the component areas associated with p . Recall that the components in \mathcal{C} are ordered by their areas, or voxel count: $\forall_{1 \leq i < k} |C_i| \leq |C_{i+1}|$. We call the largest component C_k the *background component*. The others are called *foreground components*. The importance measure ρ for a point p is defined as the total area of the foreground components:

$$\rho(p) = \frac{1}{|\partial\Omega|} \left| \bigcup_{1 \leq i < k} C_i(p) \right| \quad (8)$$

We normalize ρ by the total object surface area $|\partial\Omega|$. Note that this importance measure can only be computed for non-loop points. A regular loop-point has only one component and no foreground components (e.g. Fig. 7, point q), so that ρ is undefined. *Simplified curve skeletons* can be computed from a noisy curve skeleton \mathcal{C} by discarding all points with an importance below a certain threshold τ :

$$\mathcal{C}_\tau = \{p \in \mathcal{C} \mid \rho(p) \geq \tau\} \quad (9)$$

The threshold τ is a user-parameter which controls what is considered noise. The meaning of this parameter is intuitive: all \mathcal{C} -points representing a smaller surface area than τ are discarded. Furthermore, \mathcal{C}_τ can be considered a multiscale representation of the curve skeleton. The simplified skeletons are connected by default at all scales, because ρ is monotonic, which follows from (7). The simplified curve skeletons \mathcal{C}_τ of several shapes are shown in Figure 11, with $\tau = 10^{-3}$. Thus, curve skeleton points representing an area less than 0.1% of the object surface are discarded.

This definition of ρ has been presented in [19]. There, the curve skeleton detection was integrated with the computation of ρ . In contrast, the skeletonization method presented here computes a curve skeleton by analyzing the genus of Γ , making it much faster to compute, and making it possible to handle objects with tunnels. Furthermore, in this paper we also present a method to robustly detect junctions, as detailed in the next section.

5.4. Robust Junction Detection

As indicated in Section 4.1, the genus of $\Gamma(p)$ is a conservative criterion for detecting junctions, which is problematic if we want to use them for segmentation. Inspired by the importance measure from the previous section, we discard junctions which have small components among their foreground components. These junctions are likely the result of noise on the object boundary, or due to discretization artifacts. A point p is a junction, $p \in J_\tau$, at a certain scale τ if and only if it has at least two foreground components F_τ

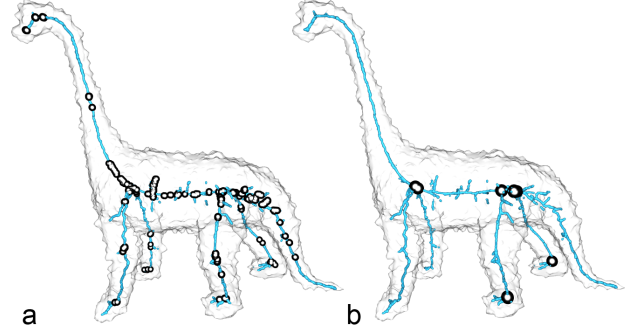


Figure 8. Conservative junctions J_0 (a). Robust junctions J_τ (b).

that are larger than τ :

$$J_\tau = \{p \in J_0 \mid |F_\tau(p)| \geq 2\} \quad (10)$$

$$F_\tau(p) = \{C_i(p) \mid 1 \leq i < k \wedge \frac{1}{|\partial\Omega|}|C_i(p)| \geq \tau\}, \quad (11)$$

where J_0 is the set of conservatively detected junctions from Section 4.1. Although related to the computation of ρ , we do not have to compute ρ to compute J_τ : We only have to compute the component sets for the conservative junctions J_0 . Equivalent to the computation of \mathcal{C}_τ , the scale τ is used to distinguish between small-scale noise and signal. For all objects we tested, a value of $\tau = 10^{-3}$ gives good results. Like the importance measure, a limitation of this method is that junctions on \mathcal{C} -loops are discarded, as these junctions have only one foreground component (Sec. 5.1).

In Figure 8a, the conservative junctions J_0 on \mathcal{C} are shown for an object with noise on its boundary, computed by analyzing the genus of Γ . Figure 8b shows the robust subset selected by applying (10).

6. Segmentation

We now present two methods for computing a shape segmentation from the critical points. By segmentation we mean the final segmentation S of the object into disjoint and connected segments S_i . In order to obtain a segmentation into meaningful components, the junctions of the curve skeleton are detected and used as critical points. Alternatively, we can pick critical points by some other (semi) automatic process. The first segmentation method (Sec. 6.1) we present is the most straightforward. It is based on the geodesic sets of the critical points, and produces a flat (non-hierarchical) segmentation at the finest scale. The second method (Sec. 6.2) is based on the component sets generated by the critical points, and produces a hierarchical, level-of-detail segmentation.

```

1:  $F \leftarrow \{C_i(p) | p \in P, 1 \leq i < k\} \cup \{C'(p) | p \in P\}$ 
2:  $F \leftarrow F \cup \{\partial\Omega\}$ 
3: procedure computeS( $\tau$ )
4:  $S \leftarrow \emptyset$ 
5: for each  $f \in F$  in ascending order of  $|f|$  do
6:   if  $\frac{1}{|\partial\Omega|} |f| \geq \tau$  then
7:      $s \leftarrow f \setminus S$ 
8:     if  $f \cap S \neq \emptyset \Rightarrow |s| > 0.1 \cdot |f \cap S|$  then
9:        $S \leftarrow S \cup \{s\}$ 
10:    end if
11:  end if
12: end for
13: end procedure
14: procedure computeLevelsOfDetail()
15: for each  $f_i, f_{i+1}$  in  $F$  in ascending order of  $|f|$  do
16:   if  $|f_{i+1}| - |f_i| > 0.1 \cdot |f_i|$  then
17:     computeS( $\frac{1}{|\partial\Omega|} |f_i|$ )
18:   end if
19: end for
20: end procedure

```

Figure 9. Algorithm for computing a level-of-detail segmentation.

6.1. Flat Segmentation

We produce a simple, flat segmentation by considering the shortest-path sets Γ of all critical points P at once. The connected components in the boundary graph due to all these shortest paths are labeled. A segmentation that is obtained in this way is shown in Figure 10b, using the shortest-path sets of the critical points shown in Fig. 10a. One property of this segmentation is that it splits tunnel-parts of the object into multiple segments.

6.2. Hierarchical Segmentation

In order to produce a hierarchical, level-of-detail segmentation, we consider the component sets C of the critical points P . In Section 5.3 we distinguished between foreground and background components. In a component set C consisting of k components $C_{1..k}$, the largest component C_k is called the background component, the remaining ones are called foreground components. The foreground components are those that one would consider meaningful and intuitively associate with the critical point. The background component is merely the remaining surface area. In Figure 6 for example, the background components are light gray. Furthermore, we observe that the foreground components combined also form a meaningful component. We denote this compound component by $C'(p) = \cup_{1 \leq i < k} C_i(p)$.

Let F be the set of meaningful components of all critical points combined. The segmentation S should be based on F , but the components in F are not disjoint. We now present an algorithm for creating a segmentation S from F

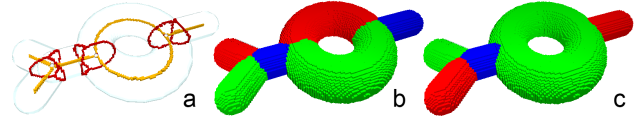


Figure 10. Critical points and their path sets (a). Flat segmentation (b). Finest level of the hierarchical segmentation (c).

consisting of disjoint segments S_i at a certain scale τ . The pseudo code of the algorithm is shown in Figure 9.

The idea is that we consider all components $f \in F$ in ascending order of their area (line 5), but only those components that are larger than the specified scale τ (line 6). The potential segment s is computed as the set difference between f and the existing segments in S : $s = f \setminus S$ (line 7). Before adding the potential segment s to S , we check whether f overlaps any existing segments from S . If not, s is added without restriction. If f overlaps, we only add s if it contributes enough to the segmentation, that is, if it adds at least 10% of the area that it overlaps (line 8). This is to prevent tiny segments due to different junctions having similar components among their component sets. This occurs for example due to the fact that junctions computed by the algorithm in Section 4 may consist of multiple voxels, having almost the same component sets. After processing all components in F , because the background components have been left out in the segmentation, the object surface might not be fully covered. Therefore, we add to F the whole boundary as the largest component (line 2), which ends up as a single segment filling up the remaining part.

In order to compute multiple segmentations from fine to coarse (line 14), we can simply consider all components from F in ascending order of area, and produce a segmentation for each of those areas. To limit the amount of generated hierarchy levels, in our implementation we only compute a hierarchy if two consecutive areas differ by at least 10% from the smaller area f_i (line 16). The different segmentations produced at the various scales actually form a hierarchy, because every segment is included in a segment from a coarser scale, by (7). Besides being hierarchical, another difference between this segmentation method and the flat one is that tunnel-parts are not segmented and are left intact (see Fig. 10c).

7. Results

We have implemented our framework in C++ and have run it on a Pentium 4, 3GHz with 1GB of RAM. In computing the shortest paths we used a cache to prevent computing the same shortest path twice. As input we used several polygonal meshes from [1, 20], voxelized using binvox [16]

for various resolutions. We used both organic and geometric objects. Figure 11 shows for four objects the simplified curve skeleton \mathcal{C}_τ and robust junctions J_τ , both with $\tau = 10^{-3}$, and three selected levels from the hierarchical segmentation. Figure 12 shows for several objects the segmentation at the finest scale. We observe that our method is able to extract fine details, such as the toes and fingers of the Armadillo and Dinopet objects, and it does not suffer from over-segmentation. The use of geodesics as borders of the surface segments has the very beneficial effect of producing sharp, non-jagged separations, since geodesics are minimally twisting curves on the surface. Additional results are available online [17].

Table 1 shows performance measurements on our framework for several objects that are shown in this paper. Columns “object”, “dim”, “ $|\Omega|$ ”, “ $|\partial\Omega|$ ” denote object name, voxel-grid dimensions, number of object voxels, and number of boundary voxels respectively. Column “ \mathcal{C} time” denotes the time in seconds to compute the non-simplified curve skeleton \mathcal{C} , including preprocessing (Sec. 4.1). Column “seg time” denotes the time for computing all levels in the hierarchical segmentation. This time strongly depends on the amount of levels generated. Column “ ρ time” denotes the time required to compute ρ on \mathcal{C} (Sec. 5.3), that is, to obtain simplified curve skeletons. Note that computing ρ is not needed for obtaining a segmentation. Column “mem” gives an indication of the peak memory usage. Computing the curve skeleton takes the most time, which can be attributed to the computation of the shortest paths and dilations. It takes relatively long for the Plane, which can be explained by the fact that the Plane has a high surface area to volume ratio, meaning that the average shortest-path between feature voxels is relatively long. The time needed to compute the robust junctions J_τ is not shown as it is negligible: up to 5 seconds for the considered objects.

8. Discussion

Following from the curve skeleton definition, the borders of the segments are piecewise geodesic. They do not necessarily follow local surface-concavities, but instead find the minimal-length path between feature points. Hence, the borders exhibit minimal twist on the surface, look natural, and are robust to noise, as shown by the Noisydino object (Fig. 12). For segmentation methods based on surface curvature, this noise could be problematic [3].

Segment borders are not based on boundary features but on the curve skeleton, capturing global geometrical (e.g. symmetrical) and topological properties. This has a number of advantages. The segmentations respect the object’s circular symmetry and are invariant for different poses of the same object. This is because the curve skeleton structure in general remains stable under deformations of the object.

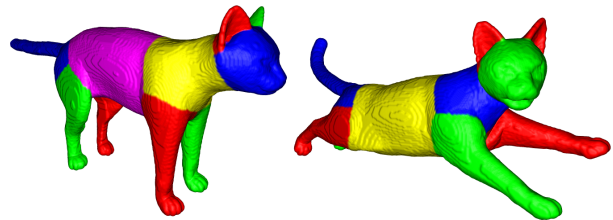


Figure 13. Pose-invariance

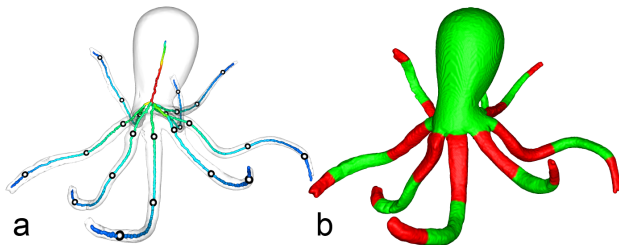


Figure 14. Manual segmentation of the Octopus.

Figure 13 shows the segmentations for two poses of the Cat (poses from [21]).

Borders are not necessarily found at curvature extrema. In the H-shape object for example, borders are found on flat parts of the surface, segmenting the tips of the H-shape. These segments are ascribed to the fact that the tips have sharp convex corners, which generate branches and junctions in the curve skeleton. A consequence of choosing only the junctions as critical points is that we do not find segment borders for curvature extrema in tubular-like parts of the object. In Figure 12 for example, the Dinopet object has no segment borders on its knees. In order to consider more the object’s geometry in addition to its topology, we could choose extra critical points on the curve skeleton at curvature extrema, computed either on the curve skeleton or original surface.

As mentioned earlier, our framework supports manual selection of critical points. In Figure 14a, we manually selected three such points on each tentacle of the Octopus. The resulting segmentation containing three segments per tentacle is shown in Fig. 14b.

9. Conclusion

We have presented a novel framework for hierarchical segmentation of 3D shapes. Being based on the formally defined curve skeleton, our framework has a solid underpinning. The produced segmentations inherit several desirable properties, such as pose-invariance, from the curve

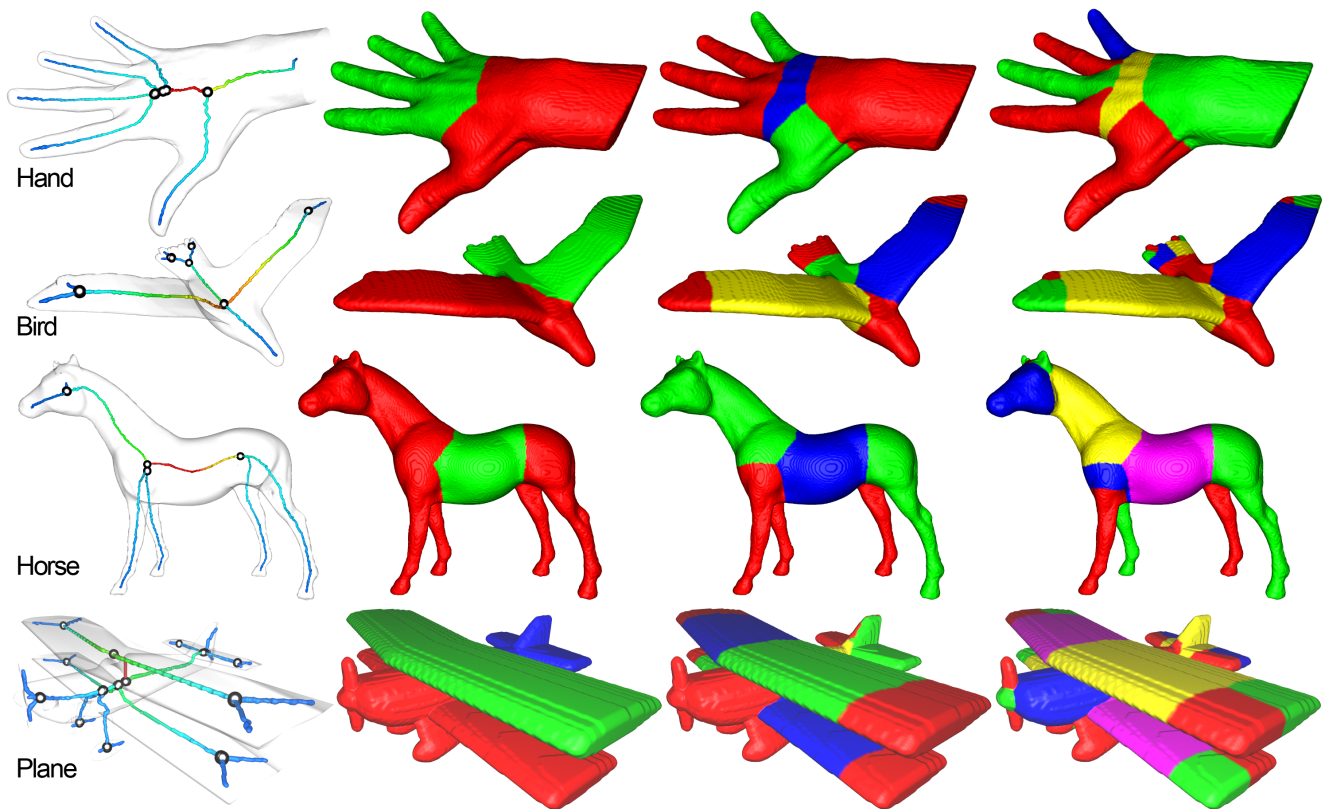


Figure 11. Simplified curve skeletons \mathcal{C}_T , and three levels of the hierarchical segmentation.

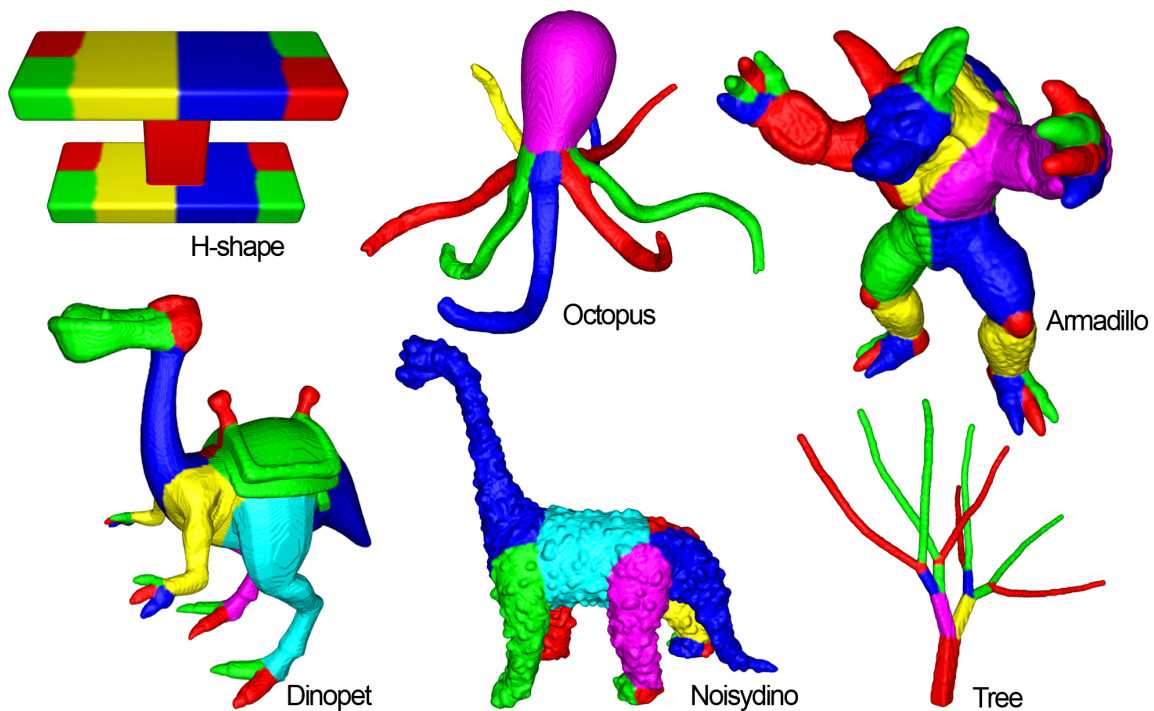


Figure 12. Segmentations at the finest scale.

Table 1. Table with measurements. See the text for details.

object	dim	$ \Omega $	$ \partial\Omega $	\mathcal{C} time (s)	seg time (s)	ρ time (s)	mem (MB)
Armadillo	188x245x207	905k	80k	42.9	41.0	7.50	447
Dinopet	334x366x180	1,810k	136k	98.3	54.2	40.0	707
Hand	366x154x257	1,300k	94k	91.8	20.5	5.57	540
Horse	366x316x171	2,038k	119k	105	33.7	10.2	660
Noisydino	125x346x365	1,421k	114k	55.2	19.4	14.0	552
Plane	217x304x98	545k	110k	228	20.5	38.5	492
Octopus	366x259x335	1,860k	154k	28.5	9.47	10.5	600

skeleton and reflect the symmetry of the object. The definition of the curve skeleton in terms of shortest geodesics gives rise to a natural skeleton-to-boundary mapping. The meaningful components are found using the curve skeleton junctions and are combined into a hierarchical, level-of-detail segmentation. Being piecewise geodesic, the segment borders are smooth and non-twisting. Interestingly, because our method is based on the curve skeleton representing the object's interior, our method produces both a surface segmentation and a corresponding volumetric segmentation. The framework supports segmentations based on the critical points defined as the curve skeleton junctions, but also defined in other ways. A voxel-based implementation is provided. Besides the segmentations, it computes multiscale curve skeletons that are robust to boundary noise, and performs robust junction detection. We showed that our framework delivers good results on both organic and geometric objects. In future work, we would like to detect additional critical points to obtain segment borders at curvature extrema of the object surface.

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